Angle-Monotone Graphs: Construction and Local Routing

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Given a point set, Construct a 'Good' Graph



















[Icking et al.'95]

July 28, 2018



Related Research



Construct a small spanning geometric graph that is

- Monotone [Mastakas and Symvonis, '17]
- Good spanner planar [Chew '89; Dobkin et al, '90; Bose et al. '10, Xia, '13]
- Good spanner nonplanar [Narasimhan and Smid (Book) 2007]
- Increasing-chord [Alamdari et al.'13; Dehkordi et al.'15; Bonichon et al.'16]

Open Question

[Dehkordi, Frati, and Gudmundsson, 2015]

Input: A point set *P*

Question: Does there exist an increasing-chord graph on P with $o(n^2)$ edges?

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This presentation

It suffices to construct an **angle-monotone graph of width 90°** with subquadratic number of edges.

Angle Monotone Path of width γ



A geometric graph in the plane is angle-monotone of width γ if every pair of vertices is connected by an angle-monotone path of width γ .

Graphs & Dilation





If $\gamma = 90^{\circ}$, then the graph is a $\sqrt{2}$ spanner

Angle-monotone graphs of width 90° are increasing chord graphs.

This presentation

- We can construct an angle-monotone graph of width 90° with O(n² log log n / log n) edges
- We can construct an angle-monotone graph of width $(90^{\circ}+\alpha)$ with $O(n/\alpha)$ edges.
- We also have a 2-local routing algorithm that can find a path of width (90°+α)

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- Add all possible edges within and between these convex polygons: O((n/log n + n²/log² n) log² n)



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- Add all possible edges within and between these convex polygons: $O((n/\log n + n^2/\log^2 n) \log^2 n)$



- Partition the point set into O(n/log n) convex polygons, each of size O(log n) ^[Urabe 1996]
- For each convex polygon, construct an angle monotone graph: O(n log n) edges in total.
- Partition each of these convex polygons into a (x, y)-convex path, a (x, -y)-convex path, a (x, y)-concave path, and a (x, -y)-concave path.

➢ For every pair of paths, we now construct an anglemonotone graph spanning these paths.

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Case Analysis :

a (x, y)-convex path a (x, y)-convex path a (x, -y)-convex path a (x, y)-convex path a (x, -y)-convex path a (x, y)-convex path a (x, y)-concave path







- An example in more general setting
 Assume that you are given a (x,y)-monotone and a
 - (x,-y)-monotone path



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An angle-monotone graph of linear size

An example in more general setting
 Assume that you are given a (x,y)-monotone and a (x,-y)-monotone path



This already covers all combinations except for when both paths are of type (x,y), or both are of type (x,-y).

An angle-monotone graph of linear size

➢ For every pair of paths, we now construct an anglemonotone graph spanning these paths.

> In each case we add at most $O(\log n \log \log n)$ edges

Case Analysis :

a (x, y)-convex path a (x, y)-convex path

> a (x, -y)-convex path a (x, y)-convex path

> > a (x, -y)-convex path a (x, y)-concave path

 $O(n^2/\log^2 n) \times O(\log n \log \log n)$

A 3-sweep graph on a point set



A 3-sweep graph on a point set



A 3-sweep graph on a point set



➢ Planar
 ➢ If θ_a = θ_b = θ_c = 60°, then the graph coincides with the well-known half-theta-6-graph.

A *k*-layer 3-sweep graph on a point set



A *k*-layer 3-sweep graph on a point set



2-Local Routing on *k*-layer 3-sweep graphs

- Send a message from the source to the destination
- The current vertex knows
 - The locations of its neighbors and their neighbors (2-hop distance).
 - The location of the destination.
 - Does not use additional memory (e.g., no routing table)

2-Local Routing on *k*-layer 3-sweep graphs



- Routing Rules
- The path taken is an angle monotone path of width (90°+α)

Open Questions

Does every point set admit an angle-monotone graph with linear number of edges?

Find the smallest γ such that every point set has a planar angle-monotone graph of width γ .

Currently it is known that $90^{\circ} < \gamma \le 120^{\circ}$.

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