

Trade-offs in Planar Polyline Drawings



Stephane Durocher



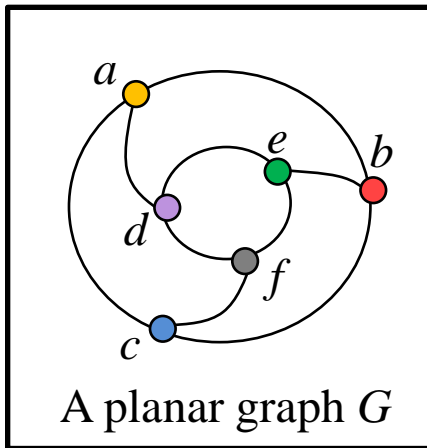
Debajyoti Mondal



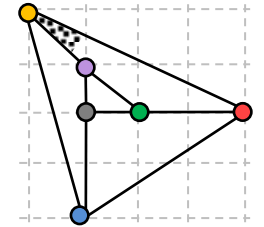
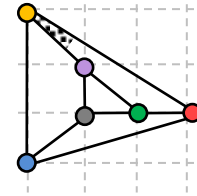
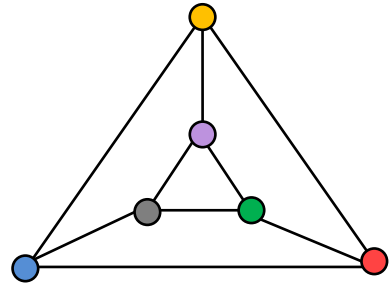
Department of Computer Science
University of Manitoba, Canada

Aesthetics of Polyline Drawings

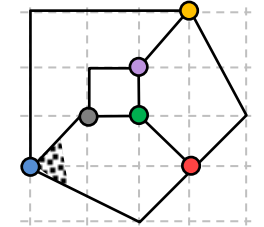
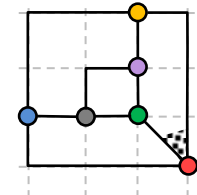
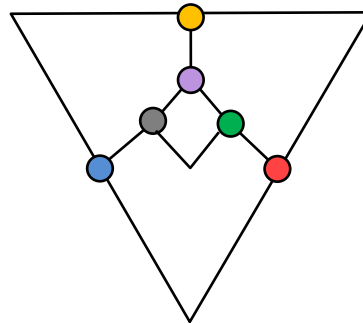
k -bend drawings of G , with angular resolution θ and area A



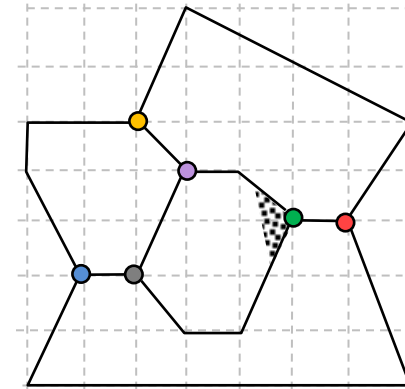
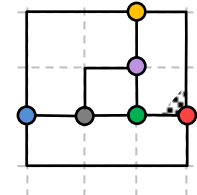
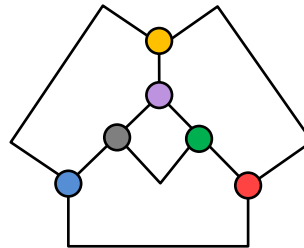
$k = 0$
 $\theta = 30^\circ$



$k = 1$
 $\theta = 90^\circ$



$k = 2$
 $\theta = 120^\circ$

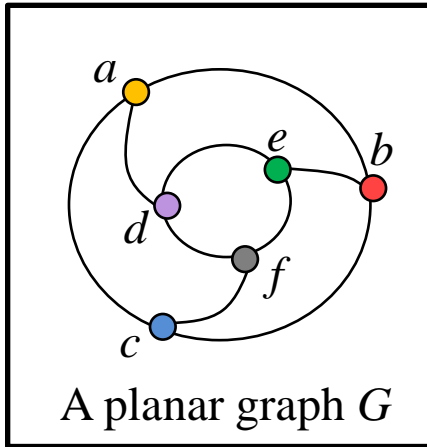


Expected Trade-offs ?

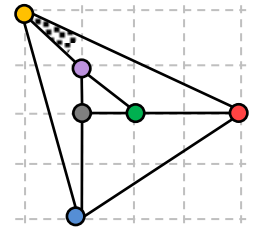
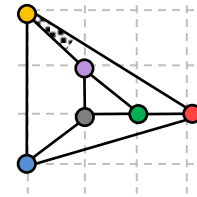
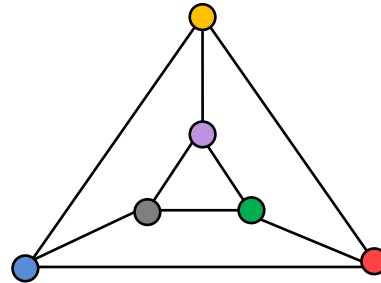
Fixed k ($\theta \uparrow, A \uparrow$)

Fixed A ($\theta \uparrow, k \uparrow$)

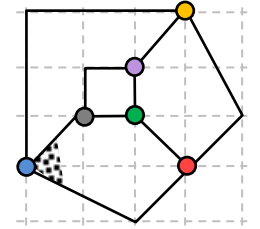
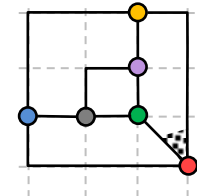
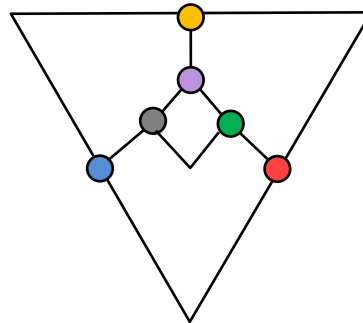
$\theta \uparrow$ ($k ? , A ?$)



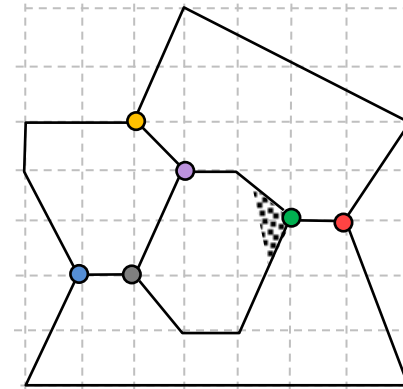
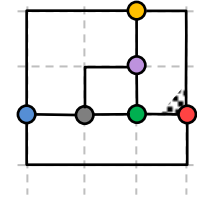
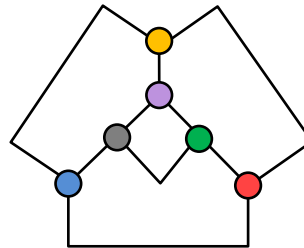
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 $\theta = 90^\circ$



$k = 2$
 $\theta = 120^\circ$

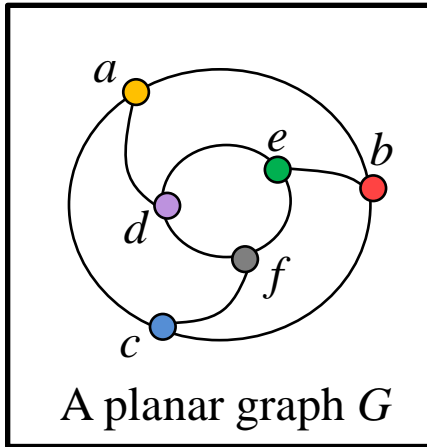


Expected Trade-offs ?

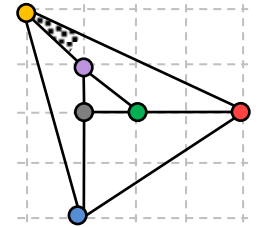
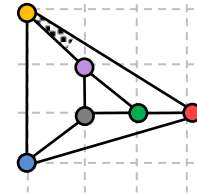
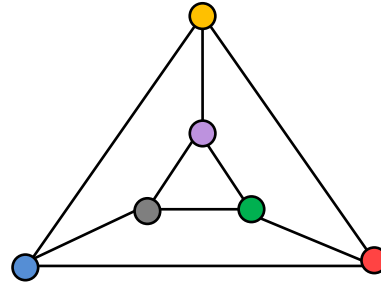
Fixed k ($\theta \uparrow, A \uparrow$)

Fixed A ($\theta \uparrow, k \uparrow$)

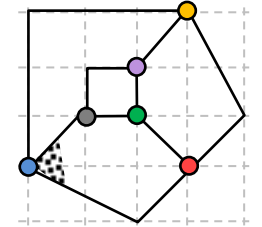
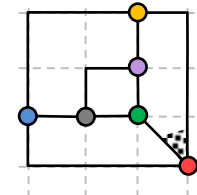
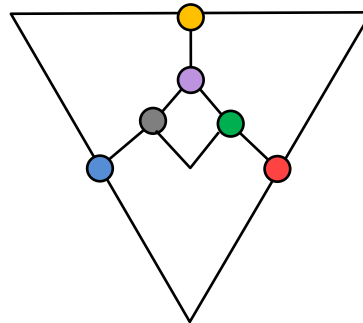
$\theta \uparrow$ ($k ? , A ?$)



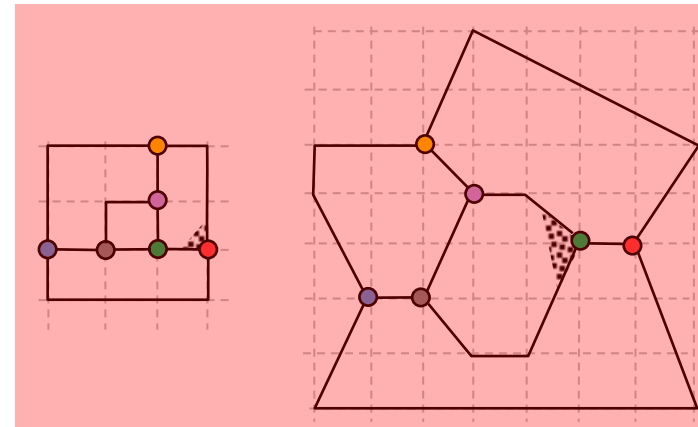
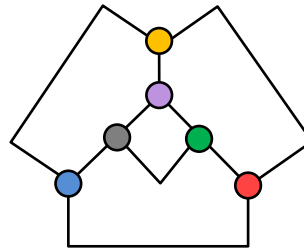
$k = 0$
 $\theta = 30^\circ$



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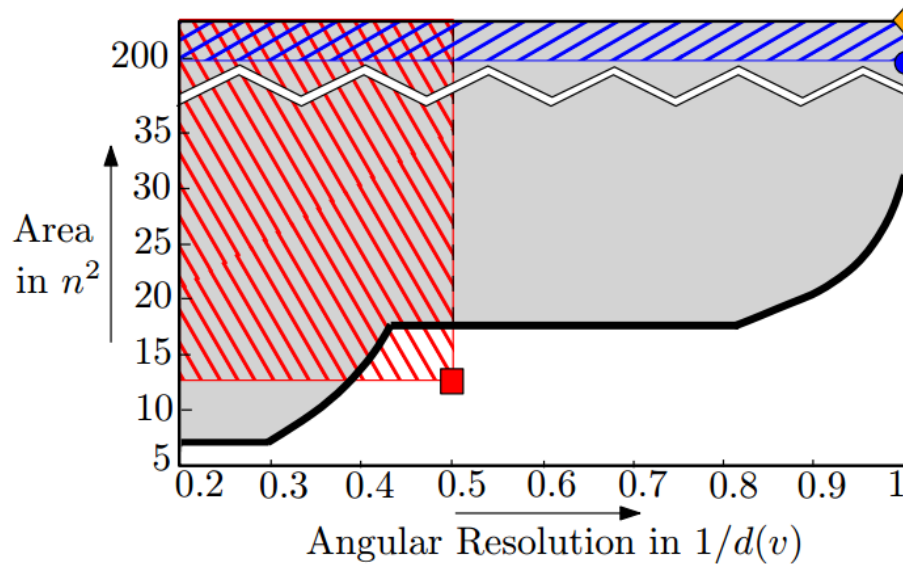
$k = 2$
 $\theta = 120^\circ$



Previous Results (on Triangulations)

<i>k</i>-Bend Drawings	Area	Angular R.	Total Bends	Reference	
$k = 0$	$0.89n^2$	$\Omega(1/n^2)$	0	Brandenburg, ENDM 2008	(A↑, θ↑)
$k = 0$	$4.50n^2$	$\Omega(1/n)$	0	Kurowski, SOFSEM 2005	
$k = 0$	$\Omega(c^{\rho n})$	$\Omega(1/\rho)$	0	Garg and Tamassia, IEEE SMC 1988	
$k = 1$	$0.45n^2$	$\Omega(1/n^2)$	$2n/3$	Zhang, Algorithmica 2010	(A↑, θ↑)
$k = 1$	$12.5n^2$	$\Omega(0.5/d_v)$	$3n$	Duncan and Kobourov, JGAA 2003	
$k = 1$	$450n^2$	$\Omega(1/d_v)$	$3n$	Cheng, Duncan, Goodrich, Kobourov, DCG 2001	
$k = 2$	$200n^2$	$\Omega(1/d_v)$	$6n$	Goodrich and Wagner, Algorithms 2000	
$k = 2$	$(6\alpha+8/3)^2 n^2$	$\frac{\alpha}{d_v(\alpha^2+1/4)}$	$5.5n$	This Presentation	
$k = 2$	$(6\beta + 2/3)^2 n^2$	$\frac{\beta}{d_v(\beta^2+1)}$	$5.5n$	This Presentation	

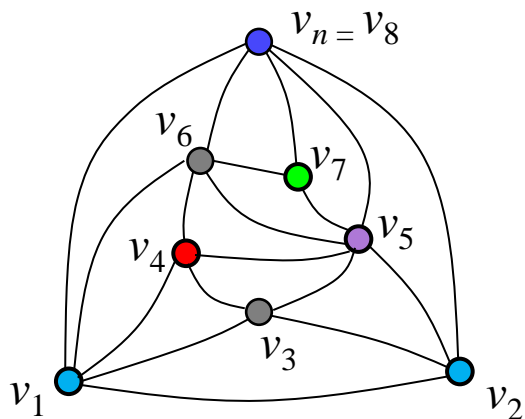
The Big Picture



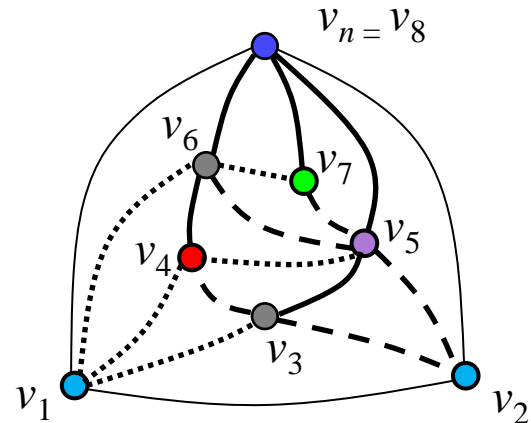
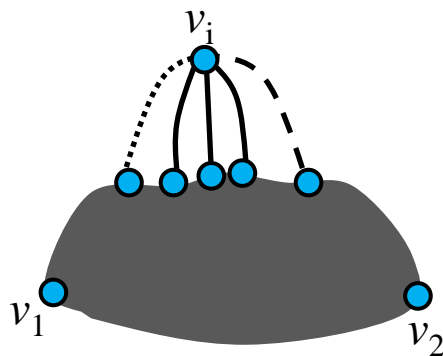
<i>k</i> -Bend Drawings	Area	Angular R.	Total Bends	Reference
■ <i>k</i> = 1	$12.5n^2$	$\Omega(0.5/d_v)$	$3n$	Duncan and Kobourov, JGAA 2003
◆ <i>k</i> = 1	$450n^2$	$\Omega(1/d_v)$	$3n$	Cheng, Duncan, Goodrich, Kobourov, DCG 2001
● <i>k</i> = 2	$200n^2$	$\Omega(1/d_v)$	$6n$	Goodrich and Wagner, Algorithms 2000
<i>k</i> = 2	$(6\alpha + 8/3)^2 n^2$	$\frac{\alpha}{d_v(\alpha^2 + 1/4)}$	$5.5n$	This Presentation
<i>k</i> = 2	$(6\beta + 2/3)^2 n^2$	$\frac{\beta}{d_v(\beta^2 + 1)}$	$5.5n$	This Presentation

Tools and Techniques

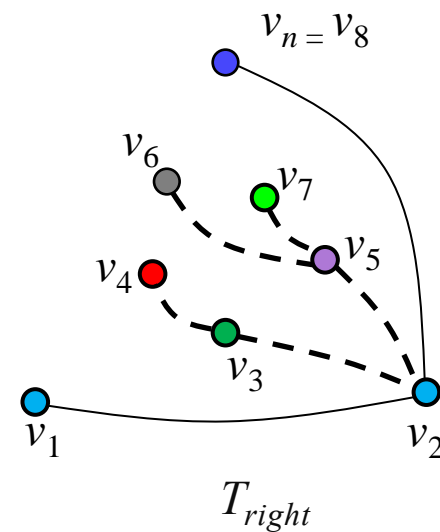
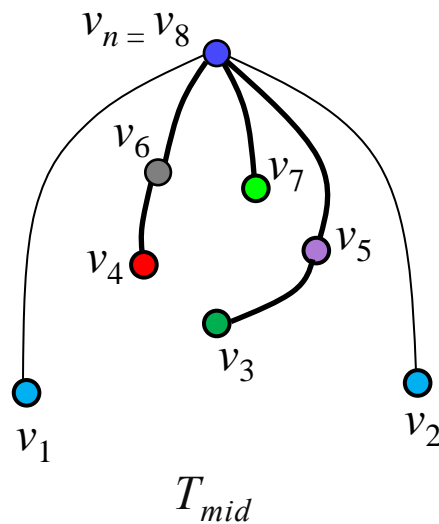
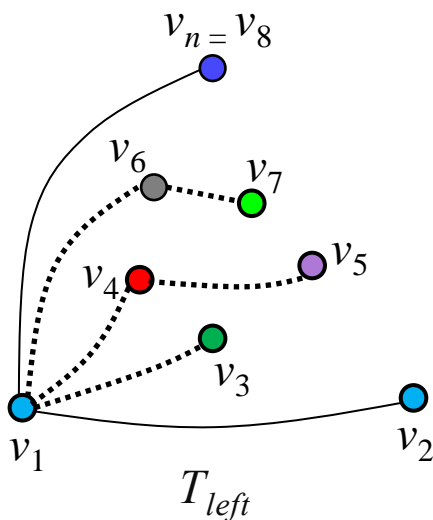
Idea: **Decomposition of a Triangulation into Trees** + *L*-Contact Representation



A Canonical Ordering of G
[De Fraysseix et al. 1988]

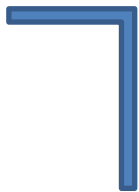
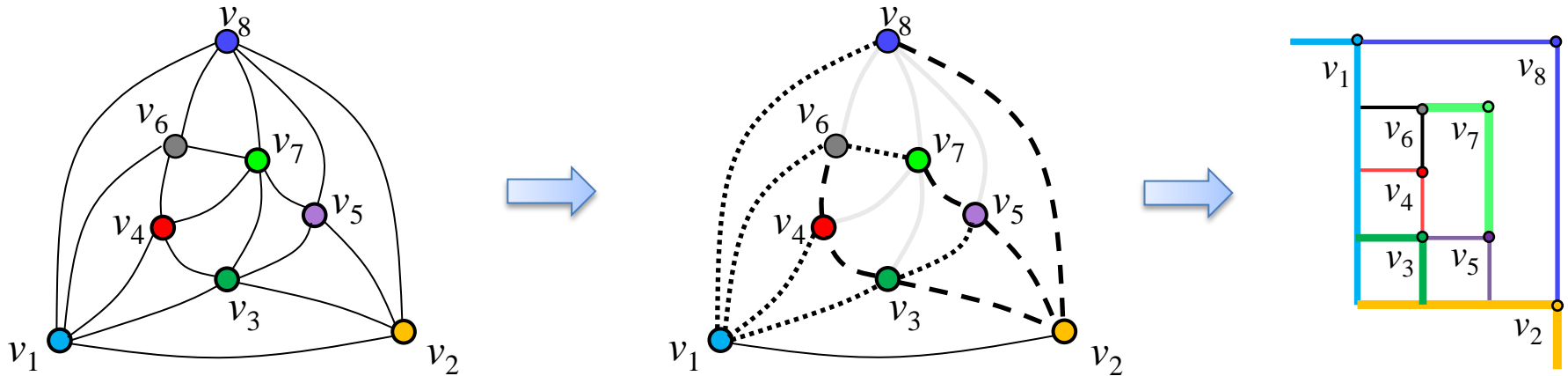


A Schnyder realization of G
[Schnyder 1990]

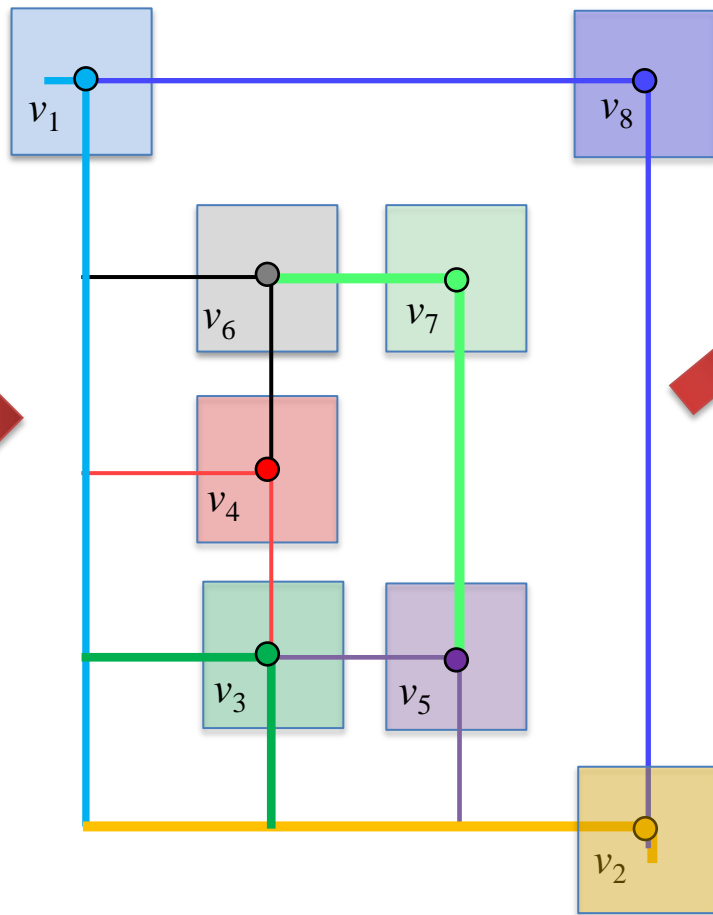
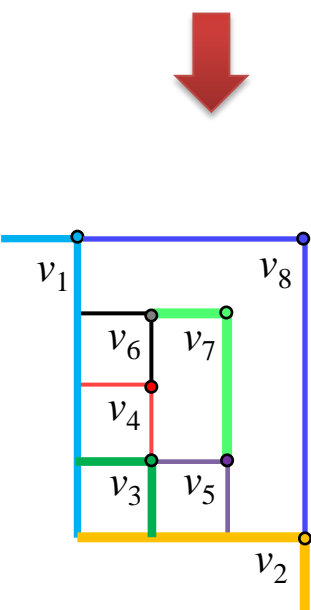
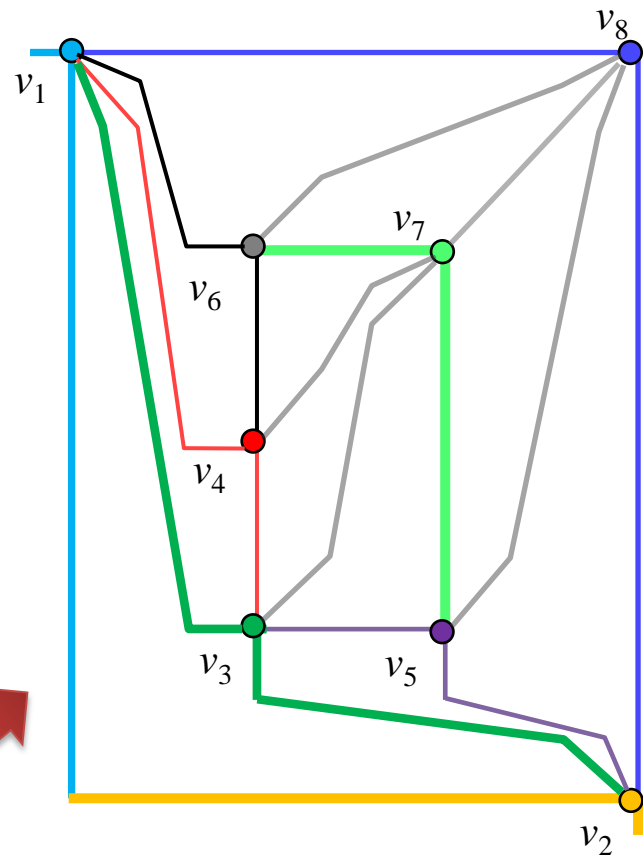
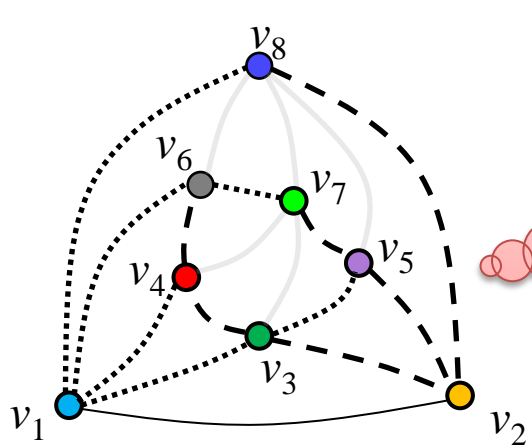


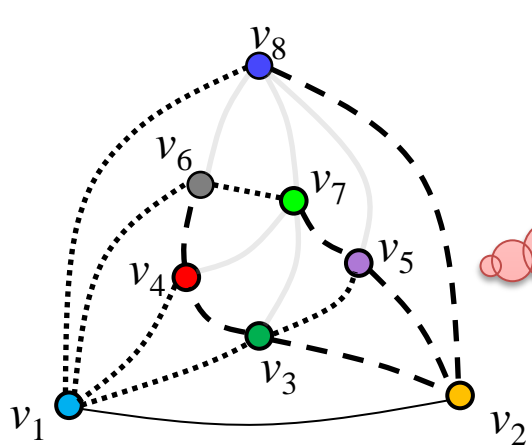
Tools and Techniques

Idea: Decomposition of a Triangulation into Trees + *L-Contact Representation*

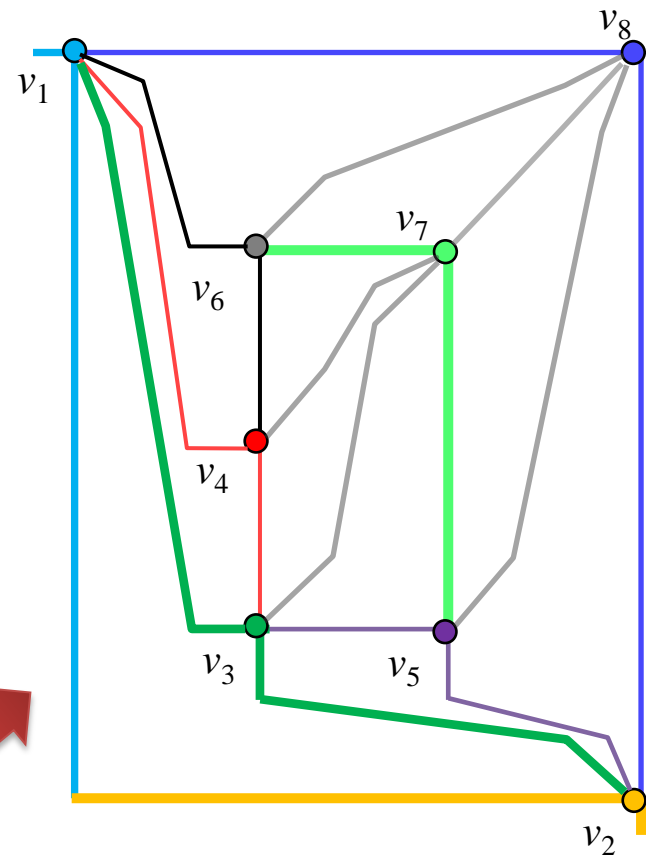
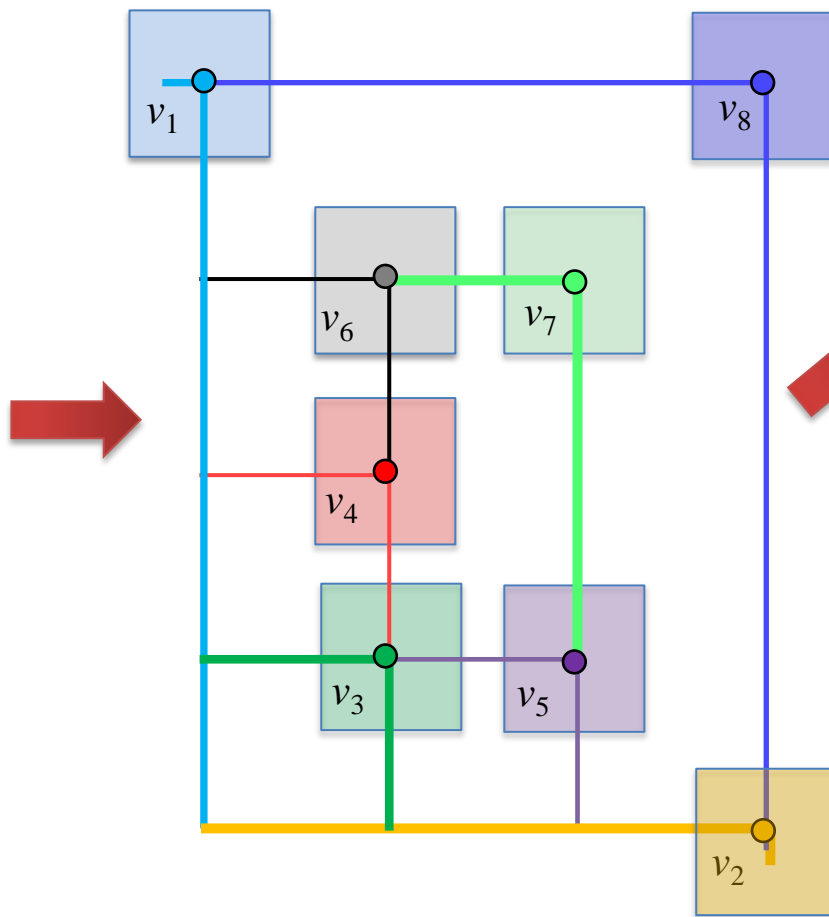
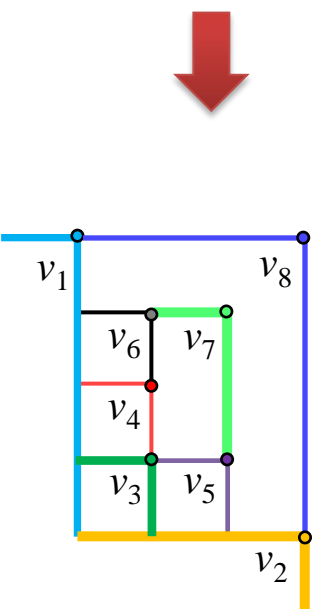


Vertices are represented with *L*-shapes
(flipped horizontally)



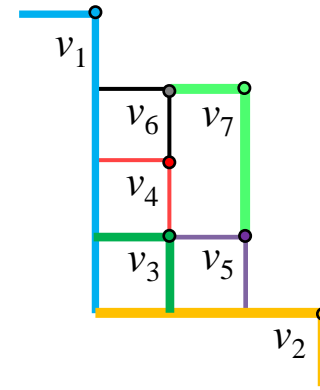
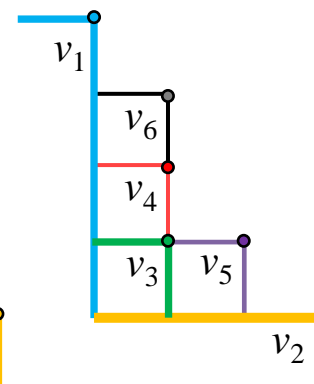
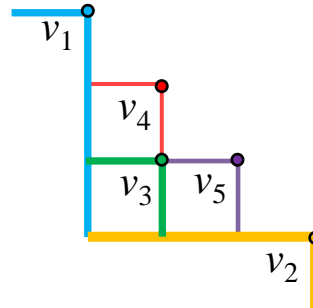
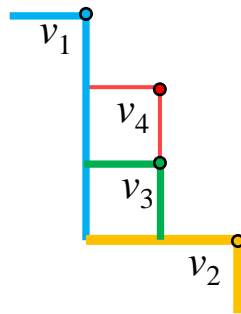
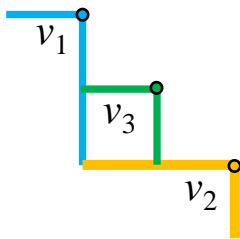
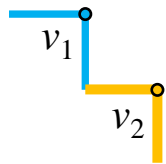
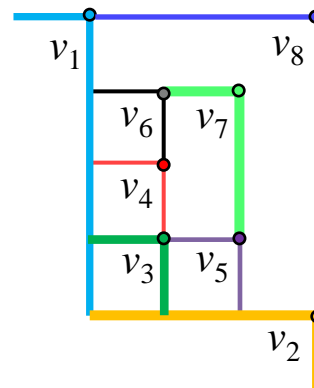
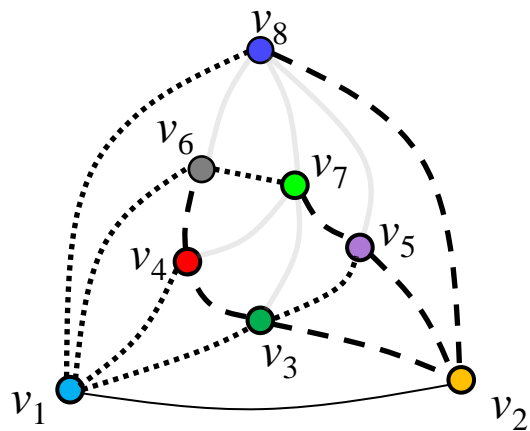


Game Plan!

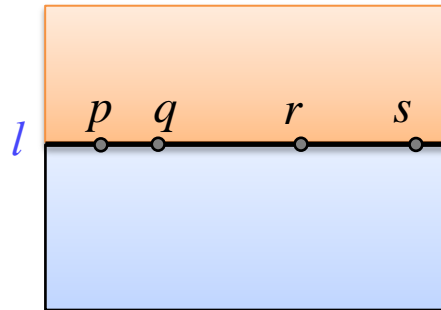
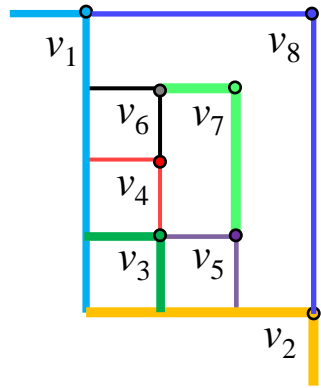


Business Meeting & Conference Dinner

Phase 1 (Contact Representation)



Phase 2 (Expansion)

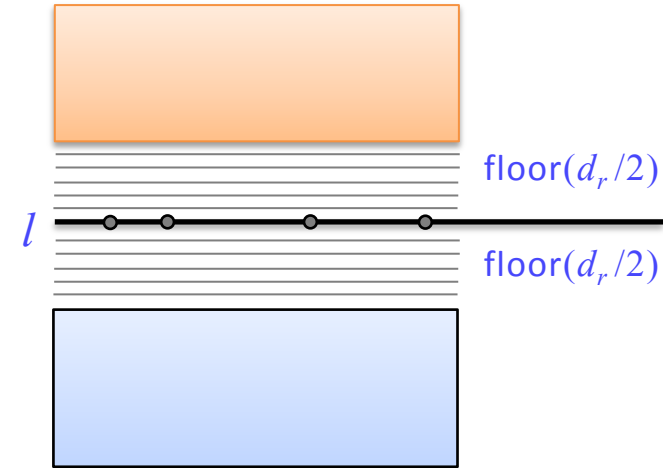


$$d_p = 6$$

$$d_q = 9$$

$$d_r = 10$$

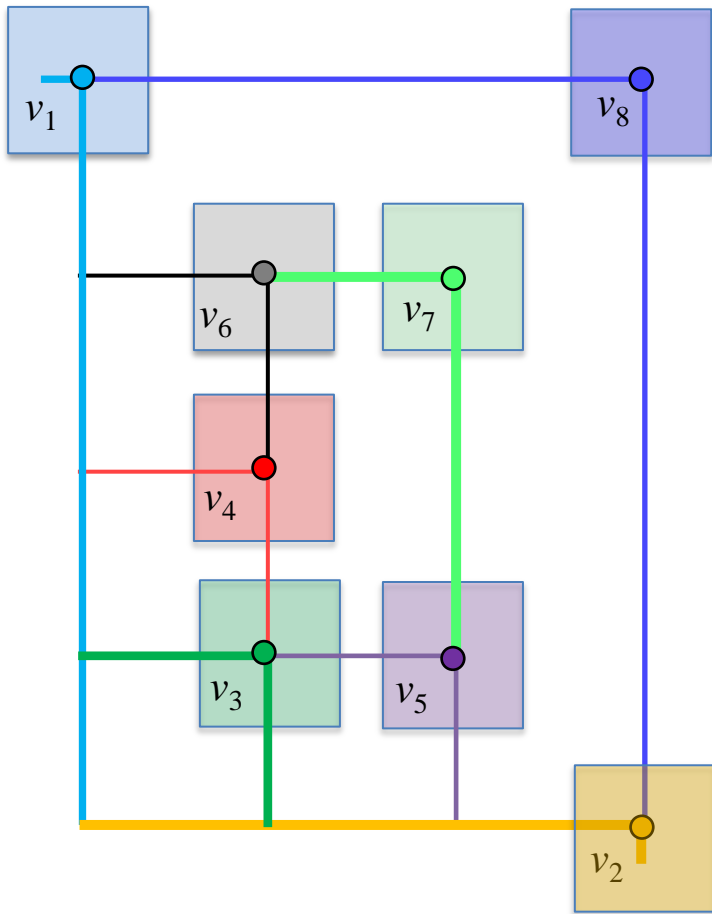
$$d_s = 4$$



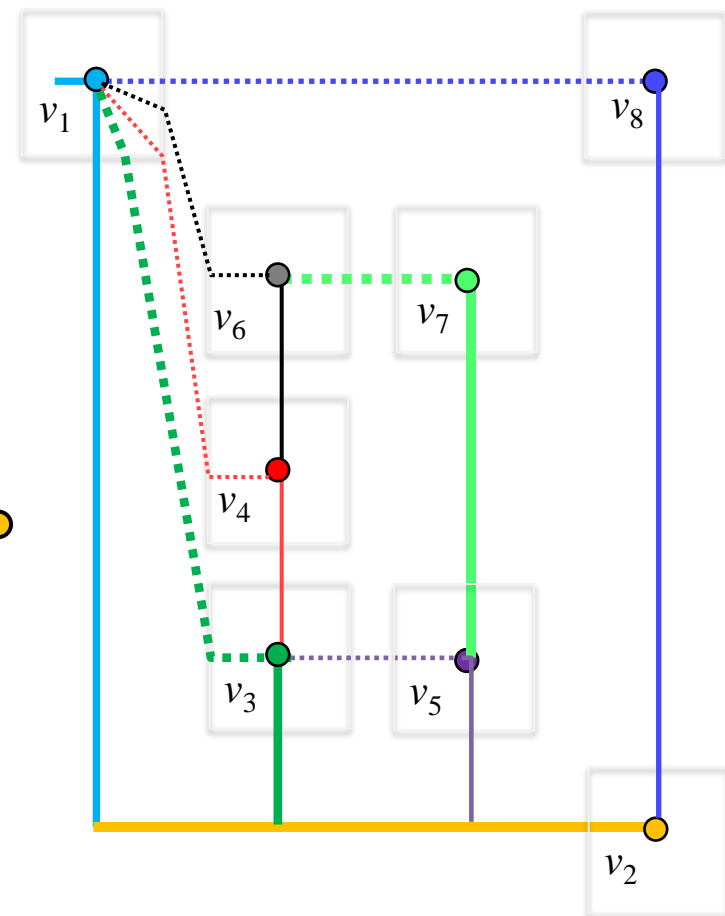
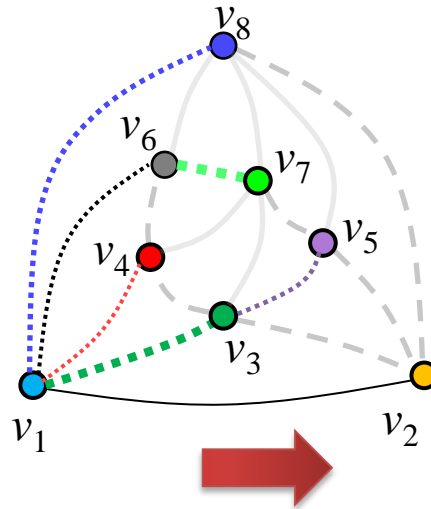
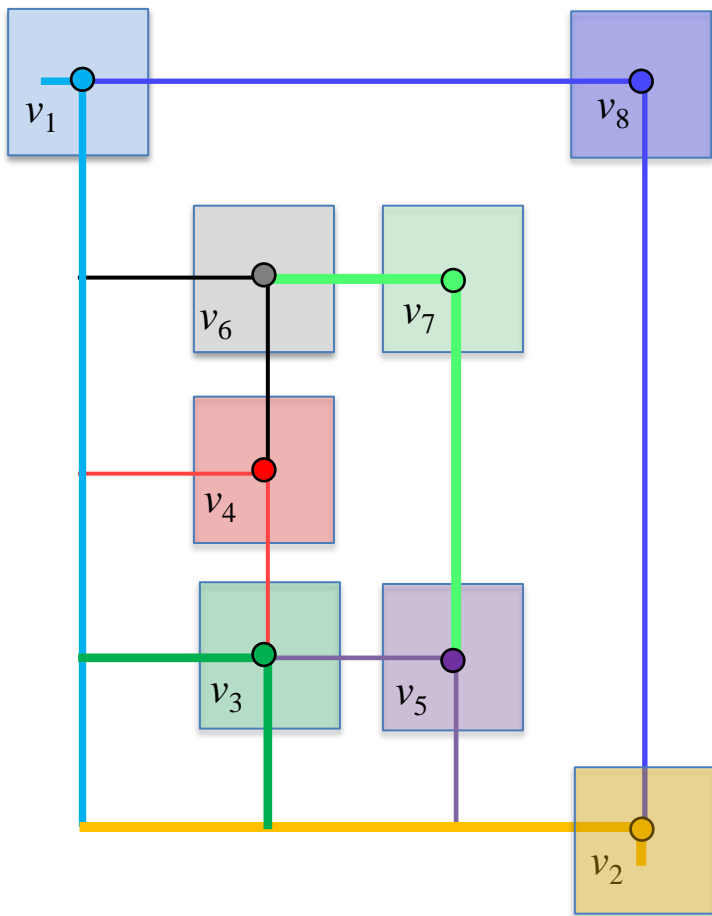
For each row l ,

- let r be the vertex on l with maximum degree.
- insert $\text{floor}(d_r/2)$ rows to each side of l .

Phase 2 (Expansion)

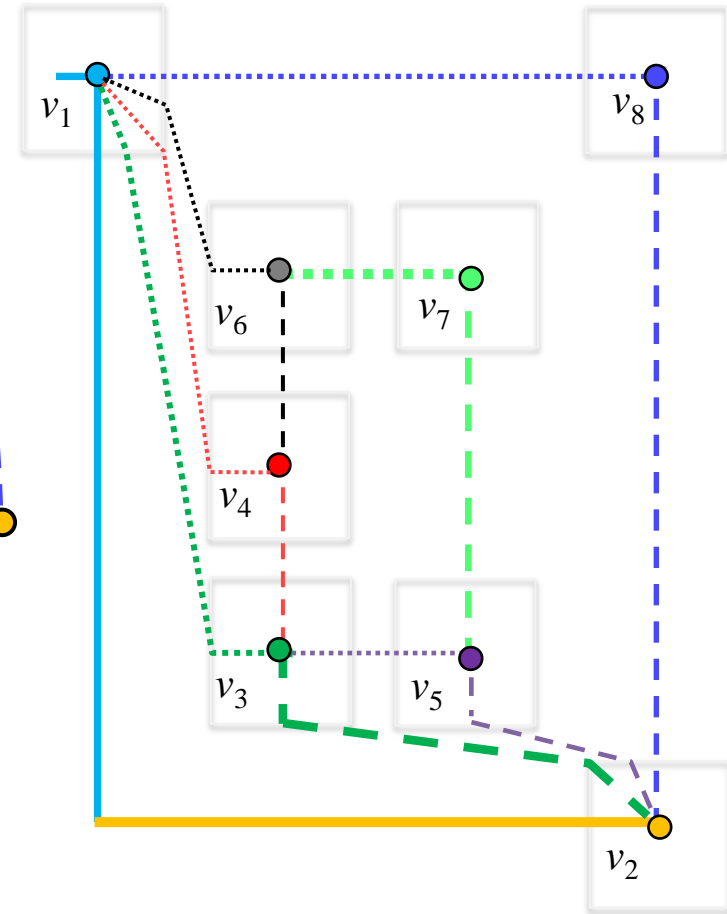
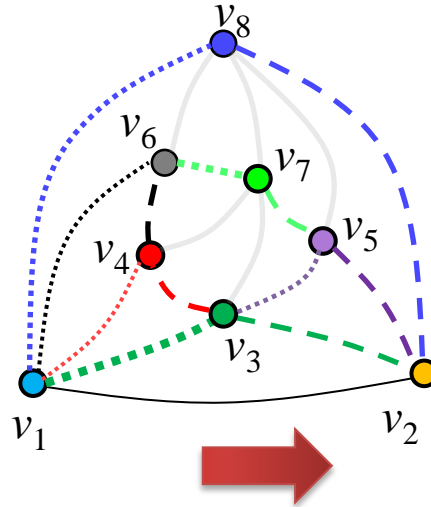
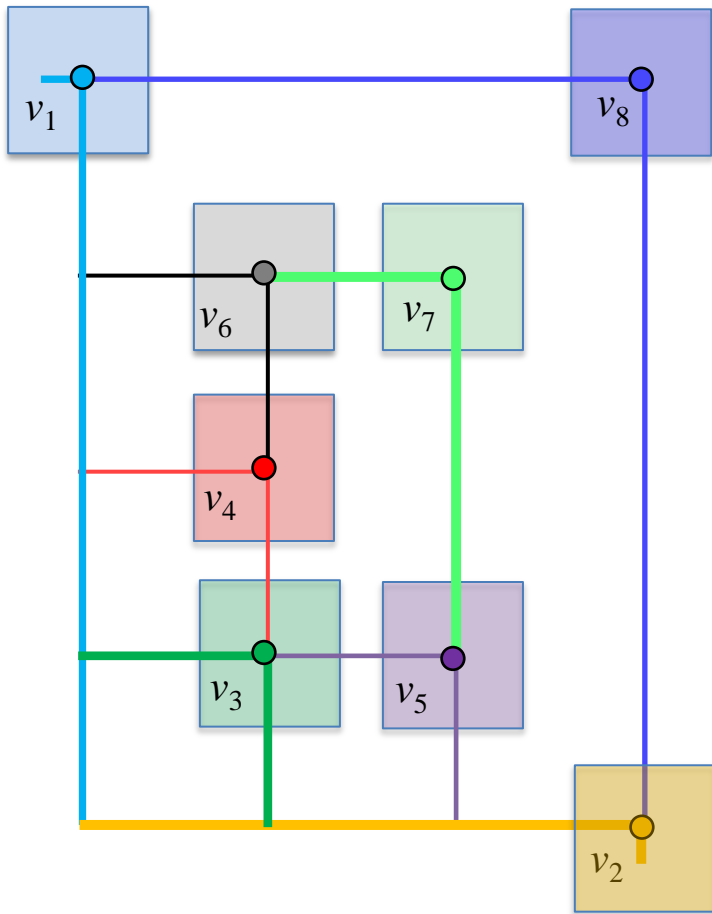


Phase 3 (Edge Routing)



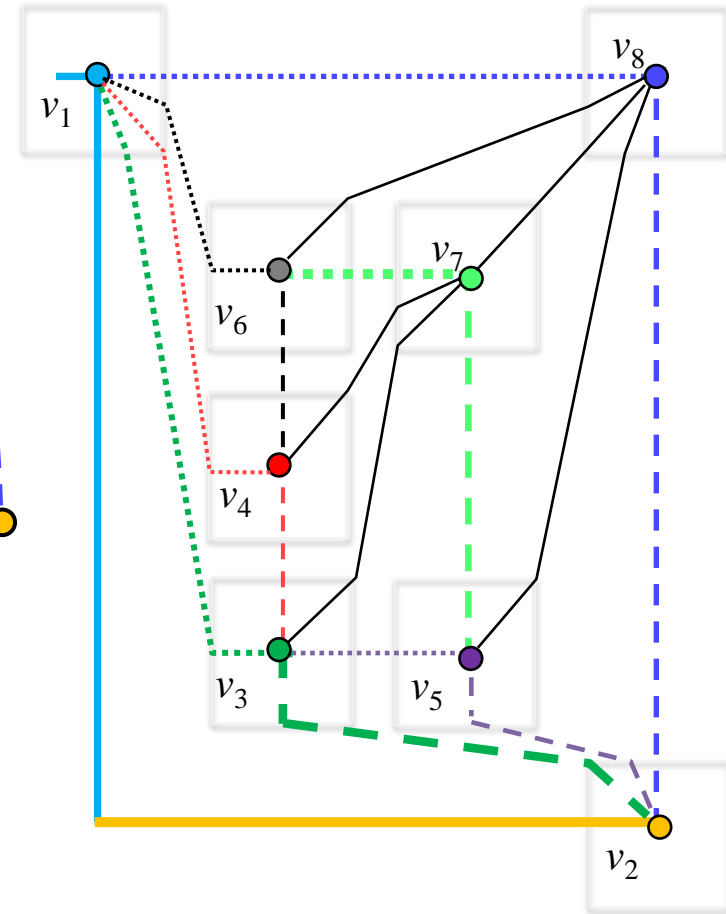
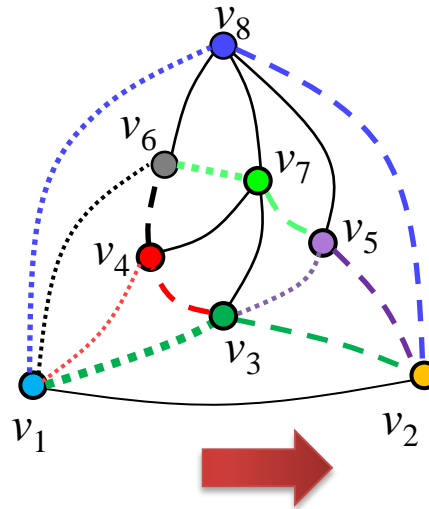
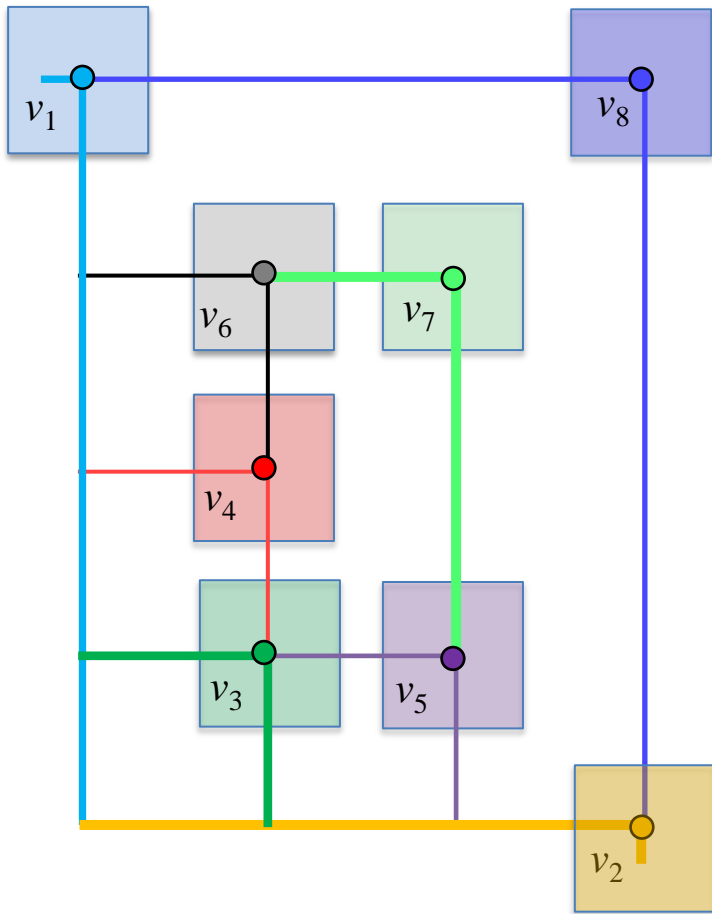
Drawing T_{left}

Phase 3 (Edge Routing)



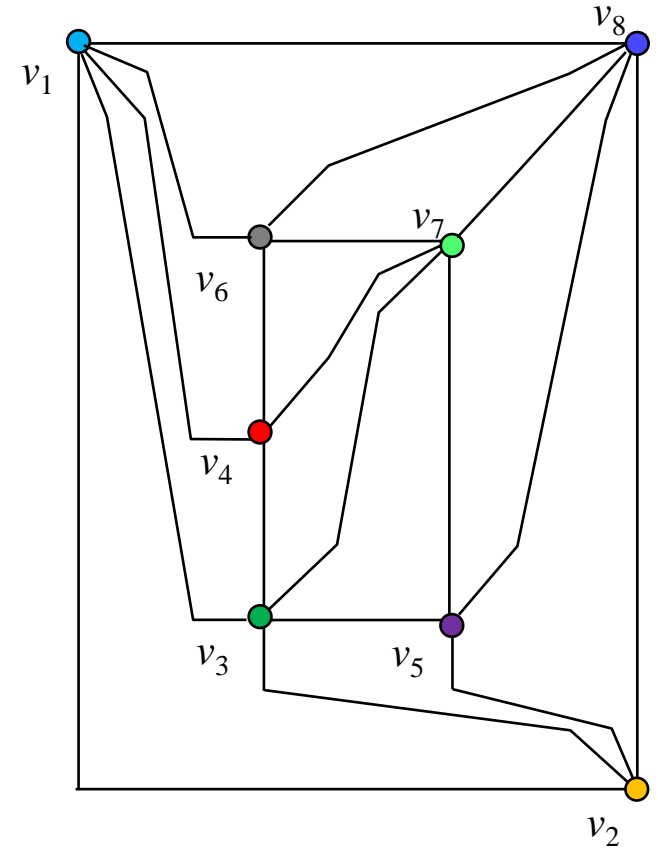
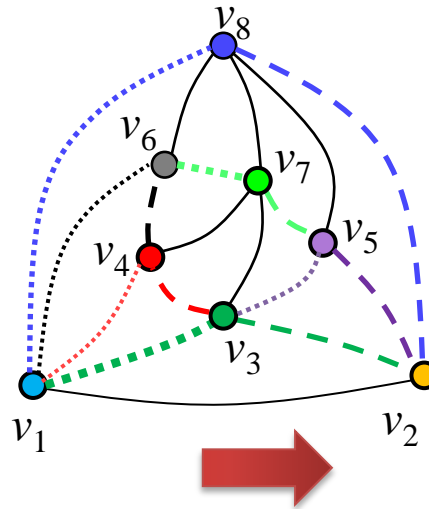
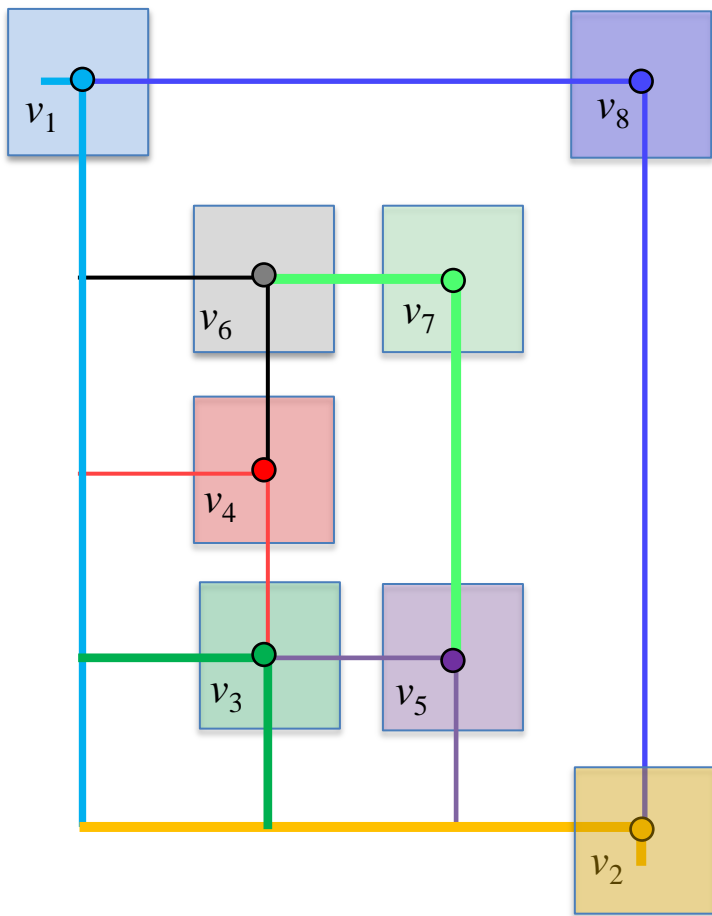
Drawing T_{right}

Phase 3 (Edge Routing)



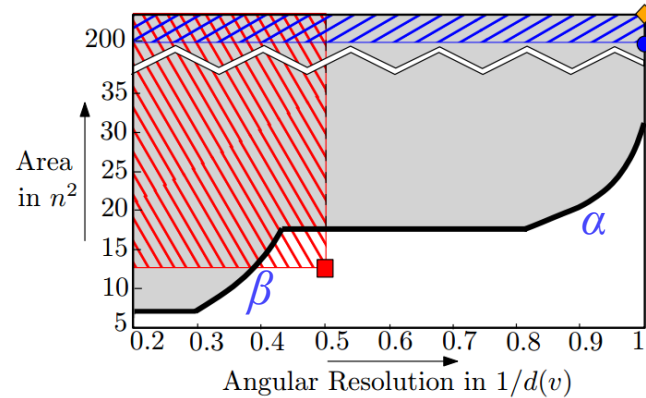
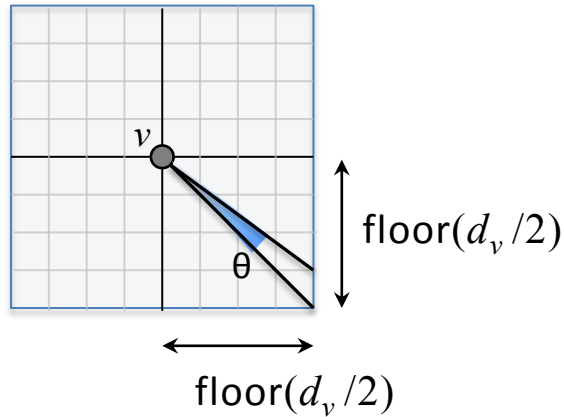
Drawing T_{mid}

Phase 3 (Edge Routing)

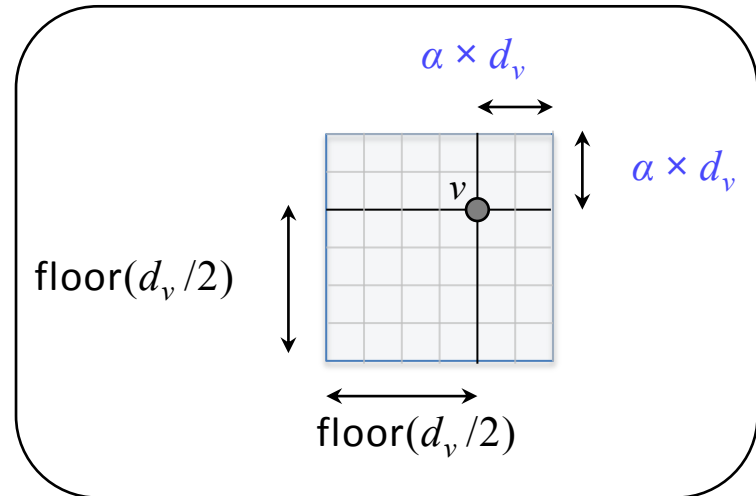
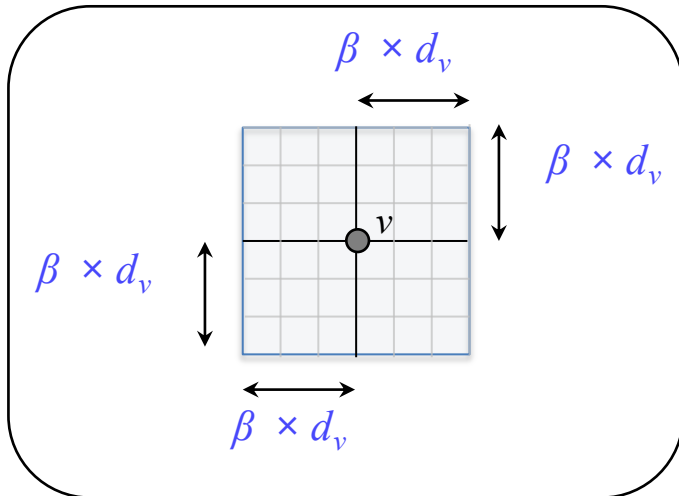


Final Output

Where does the trade-off come from?



... Because we have other options 😊



Area	Resolution
$(6\beta + 2/3)^2 n^2$	$\frac{\beta}{d_v(\beta^2+1)}$

Area	Resolution
$(6\alpha + 8/3)^2 n^2$	$\frac{\alpha}{d_v(\alpha^2+1/4)}$

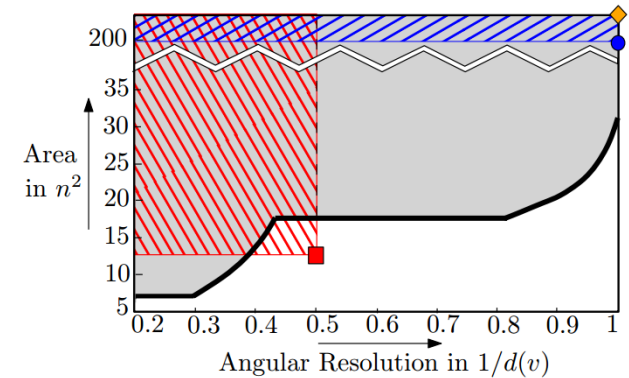
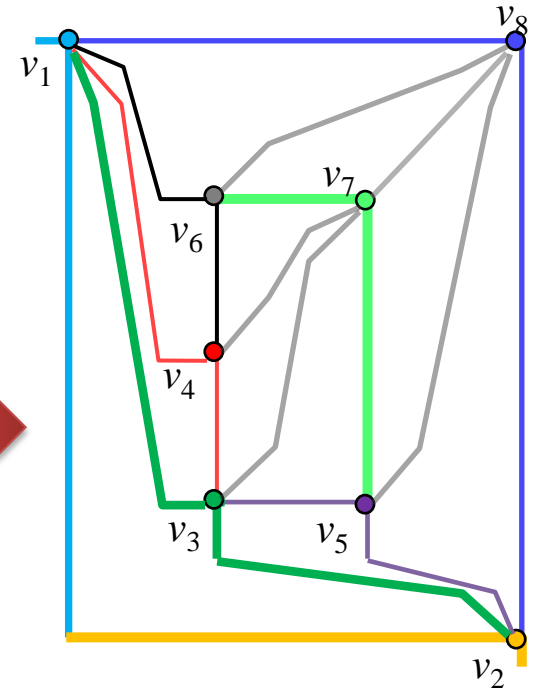
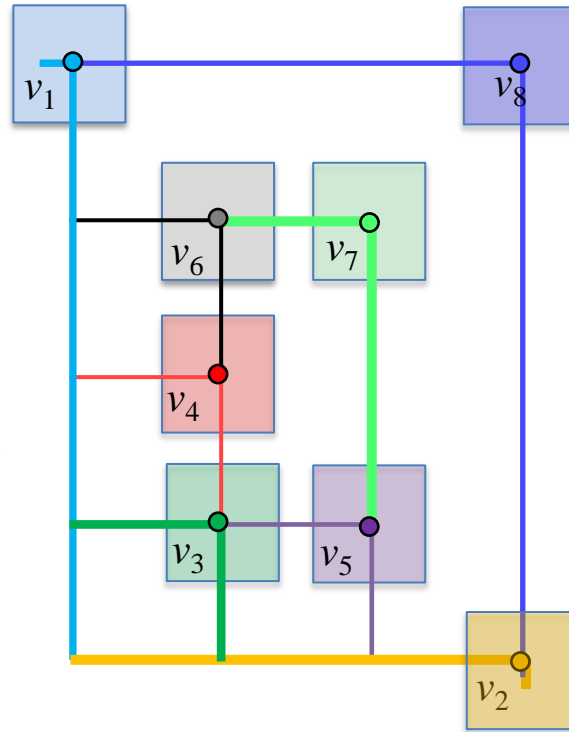
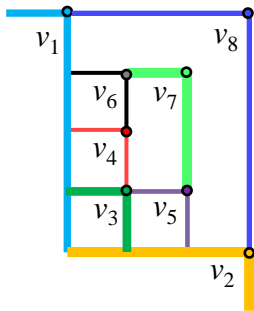
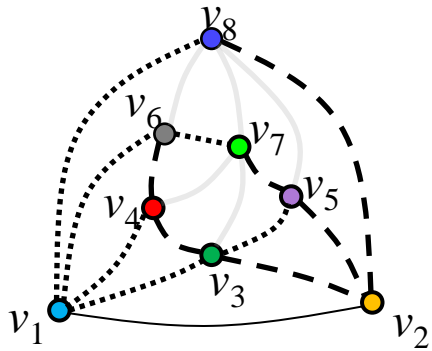
Future Research

Better Trade-offs: Can we achieve better trade-offs between area and angular resolution for 2-bend polyline drawings? Find such trade-offs for 1-bend and 0-bend drawings.

Generalization: Examine trade-offs between area and angular resolution for more general graphs such as 1-planar graphs and thickness 2-graphs.

Different Aesthetics: Examine smooth trade-offs among different aesthetic criteria for other styles of drawing graphs?

Thank You..



Better Trade-offs?

Generalization?

Different Aesthetics and Styles?